



## A Modified Method for Optimizing Area and Speed in FFT Processors

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**Abstract:**—Proposed work is a new approach to design FFT algorithm, for that significant logics are made on the basic butterfly and total numbers of multiplications are reduced. Conventionally, twelve basic butterflies are required for DIF-FFT, but in this work only four basic butterflies are used to perform the same work that is done by twelve butterflies in the conventional one. Apart from this, total four multiplications, one addition and one subtraction are required for performing multiplication between two complex numbers. But in proposed work these four multiplications are reduced into only three multiplications on the cost of one extra addition and two extra subtractions, still it is a fair deal if area and speed are concerned. After the reduction in total number of multipliers of FFT, the rest of the multiplications are done by Vedic multiplication technique. Here Vedic multiplication technique is used for multiplication part because it is observed that Vedic methodology is much better as far as speed is concerned. Thus proposed work is a two layer optimization approach which makes FFT area and speed optimized.

### 1. INTRODUCTION

Electronic signal is an electric current used to convey data from one location to another. Signals are of two types namely- Discrete signals and Continuous signals.

Discrete signal number of elements in the set, as well as the possible values of each element are finite, are countable and can be represented with computer bits.

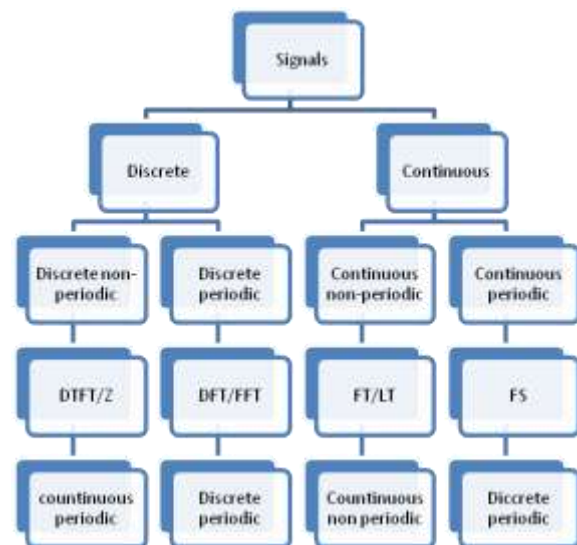


Figure 1: Hierarchy of Fourier Transform

The algorithm used to process discrete as well as continuous signals is called Fourier transform. But based on the type of signals they process, this Fourier transform is further classified into different categories namely-

1. DTFT or Discrete Time Fourier Transform.
2. DFT or Discrete Fourier Transform.
3. FT or simple Fourier Transform
4. Fourier series.

Since DFT handles a finite amount of data it and it can be implemented in computers by numerical algorithms. These implementations usually employ efficient FFT (Fast Fourier Transform) algorithm. FFT is one of the most powerful tools in digital signal processing applications and it is also the basic transformation employed by the latest wireless communication standards. Fast Fourier Transform (FFT) processor is widely used in different applications, such as-WLAN, Image process, Spectrum measurements, Radar and multimedia communication services.

## 2. MOTIVATION

In many real-time DSP applications, speed is the prime target and achieving this may be done at the expense of the accuracy of the arithmetic operations. Signal processing deals with signals distorted with the noise caused by non-ideal sensors, quantization processes, amplifiers, etc., as well as algorithms based on certain assumptions, so inaccurate results are inevitable.

There has been extensive work on high speed multipliers at technology, physical, circuit and logic levels. A system's performance can be measured by the working of the multiplier because the multiplier is normally the slowest element in the any system. Also, it is generally the most area consuming. So, optimizing the area and speed and power of the multiplier is a major design issue.

## 3. PROBLEM STATEMENT AND ENCOUNTER

FFT is largely used in many applications but whatever done till now in the field of FFT and Vedic mathematics was although an achievement in itself but was not sufficient and perfect since it consists of few limitations.

After going through reference papers it is observed that-

- Reference paper 1, they have utilized the elements of target device only

and did not use any additional multiplier because of which their area was although reduced but they could do nothing for speed enhancement.

- Reference paper 2, they had used pipelining architecture to improve the throughput, but along with this additional multipliers are also used that had negative impact on area consumption.
- Reference paper 3, they had used Urdhva Triyagbhyam method to perform multiplications. And the carry bits generated after multiplication are added by using tree addition structure. If only speed is concerned then tree addition structure is the best method but if area is concerned then there are other methods which are even better than tree addition structure.

The problem encountered in these base papers is tried to be solved in the proposed work by the following way-

- Total number of basic butterflies of DIF-FFT is reduced to  $1/3^{\text{rd}}$  so that area consumption could be reduced.
- Total number of multipliers required for complex multiplication is reduced by new proposed method.
- Vedic multiplier (Urdhva Triyagbhyam) is used to increase the speed.

## 4. PROPOSED TECHNIQUE

Consider two inputs  $Z1=(x1+iy1)$  and  $Z2=(x2 + iy2)$ . On taking conventional multiplication of these inputs they gives the outputs as- Real Part(R) =  $(x1x2-y1y2)$  and Imaginary Part (I) =  $(x1y2+y1x2)$ . On implementing it requires four multipliers and one adder and one subtractor.

But if the two inputs are multiplied by the proposed approach the outputs are given as

$$\text{Real Part(R)} = x_1(x_2+y_2)-y_2(x_1+y_1) \text{ \& Imaginary Part(I) = } x_1(x_2+y_2)-x_2(x_1-y_1)$$

Upon adding the above two terms (R and I) it gives the same value as simple multiplication. But implementation of R and I requires three multipliers and two adder and three subtractors (term  $x_1(x_1+x_2)$  is counted once because it is repeating in real and imaginary part), so one multiplier is reduced on cost of one adder and two subtractor.

Proposed complex multiplication need one extra adder and two extra subtractors on the cost of one reduced multiplier.

A 16 bit adder need 16 Full adder and 16 bit subtractor need 16 Full adders with 16 XOR gates. But one 16 bit multiplier needs  $16 \times 16 = 256$  AND gate and  $32 \times 15 = 480$  Full adder (for conventional multiplication) and this can be reducing maximum up-to 75% of conventional requirement even if advance multiplication techniques like Wallace, Vedic, booth etc. are used.

Therefore it can be concluded that still one adder and two subtractions is a better deal instead of using one 16 bit multiplier.

**Table 1: Comparison of no. of add/sub and multiplication between conventional and proposed work.**

64 Points FFT			
Conventional		Proposed	
Multiply	Add/Sub	Multi- ply	Add/ Sub
$6 \times 32 \times 4 = 768$	$6 \times 32 \times 2 = 384$	$6 \times 32 \times 3 = 576$	$6 \times 32 \times 5 = 960$
8 Points FFT			
Conventional		Proposed	
Multiply	Add/Sub	Multi- ply	Add/ Sub
$3 \times 4 \times 4 = 48$	$3 \times 4 \times 2 = 24$	$3 \times 4 \times 3 = 36$	$3 \times 4 \times 5 = 60$

Let us take an example to explain the above discussed method more clearly-

Suppose  $z_1 = 3.25 + 3j$  and  $z_2 = 7.5 + 1.17j$  are two inputs to be multiplied. The real and imaginary parts after their multiplication are found out as-

$$\begin{aligned} R &\Rightarrow 3.25(7.5+1.17)-1.17(3.25+3) \\ &\Rightarrow 3.25(8.67)-1.17(3.25) \\ &\Rightarrow 28.1775-3.8025 \\ &\Rightarrow 20.865 \end{aligned}$$

$$\begin{aligned} I &\Rightarrow 3.25(7.5+1.17)-7.5(3.25-3) \\ &\Rightarrow 3.25(8.67)-7.5(.25) \\ &\Rightarrow 28.1775-1.875 \\ &\Rightarrow 26.3025 \end{aligned}$$

Let's have the above example in binary form

$$X_1(3.25) \Rightarrow 00000000011.0100 \text{ and}$$

$$X_2(7.5) \Rightarrow 00000000111.1000$$

$$Y_1(3) \Rightarrow 00000000011.0000 \text{ and}$$

$$Y_2(1.875) \Rightarrow 00000000001.0011$$

$$X_2+Y_2(8.6875) \Rightarrow 00000001000.1011$$

$$X_1+Y_1(6.25) \Rightarrow 00000000110.0100$$

$$X_1-Y_1(0.25) \Rightarrow 00000000000.0100$$

$$\{ X_1 * ( X_2 + Y_2 ) \} ( 28.234375 ) \Rightarrow 00011100.00111100$$

$$\{ Y_2 * ( X_1 + Y_1 ) \} ( 7.421875 ) \Rightarrow 00000111.01101100$$

$$\{ X_2 * ( X_1 - Y_1 ) \} ( 1.875 ) \Rightarrow 00000001.11100000$$

$$\{ X_1 * ( X_2 + Y_2 ) - Y_2 * ( X_1 + Y_1 ) \} ( 20.8125 ) \Rightarrow 00010100.11010000$$

$$\{ X_1 * ( X_2 + Y_2 ) - X_2 * ( X_1 - Y_1 ) \} ( 26.359375 ) \Rightarrow 00011010.01011100$$

So the final Real part is  $R = 20.8125$  and imaginary part is  $I = 26.359375$ .

## 5. COMPARATIVE RESULTS

**Table 2: Comparison table of Base 1 and Proposed work (8 point FFT)**

	Ref 1	Proposed
No. of slices	1989	1927
No. of 4 input LUTs	3627	3405
Logical delay	600ns	32.557ns
No. of multipliers	Nil	Nil

Since the Ref paper 2 is the work done on 4 point FFT hence the table shown below is the comparative results of 4 point FFT

**Table 3: Comparison table of Base 2 and proposed work (4 point FFT)**

	Ref 2	Proposed
No. of slices	562	637
No. of LUTs	-	1127
Logical delay	31.55ns	23.780ns
No. of multiplier	12(9x9 multipliers)	NIL

Here in the proposed work, no. of slices are more as compared to Ref 2 but it can be seen that Ref 2 have used additional multipliers which is nil in proposed case. Therefore the overall area of the proposed work is still less as compared to Ref 2.

## 6. CONCLUSION

The proposed work is a double layer optimization technique where firstly the reduction in total number of basic butterflies along with reduction in total number of multipliers of FFT makes the proposed FFT, area efficient and secondly the number of multiplications that were left is done via Vedic multiplication approach which increases the speed.

Finally the synthesized and simulated results so obtained show that the proposed design is producing the expected and efficient results as compared to Reference papers.

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